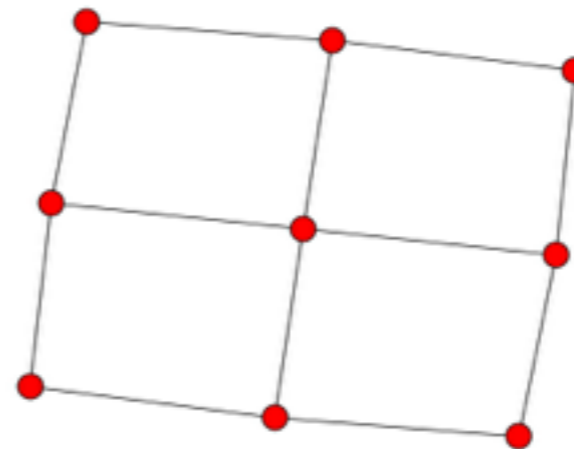
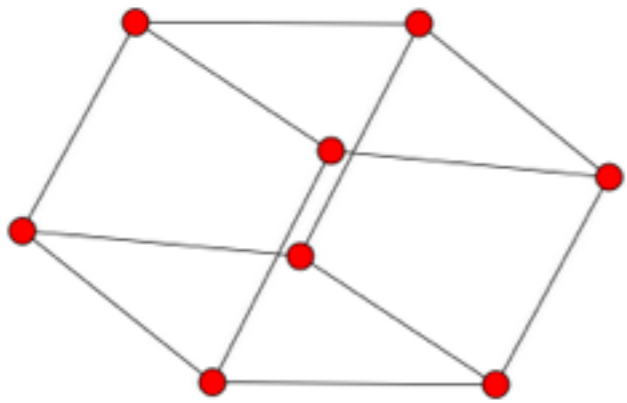
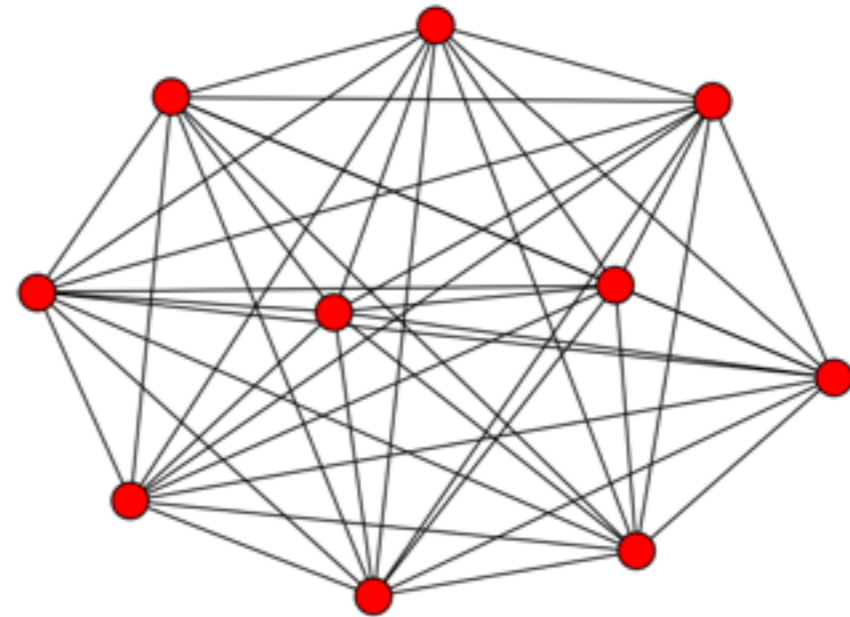


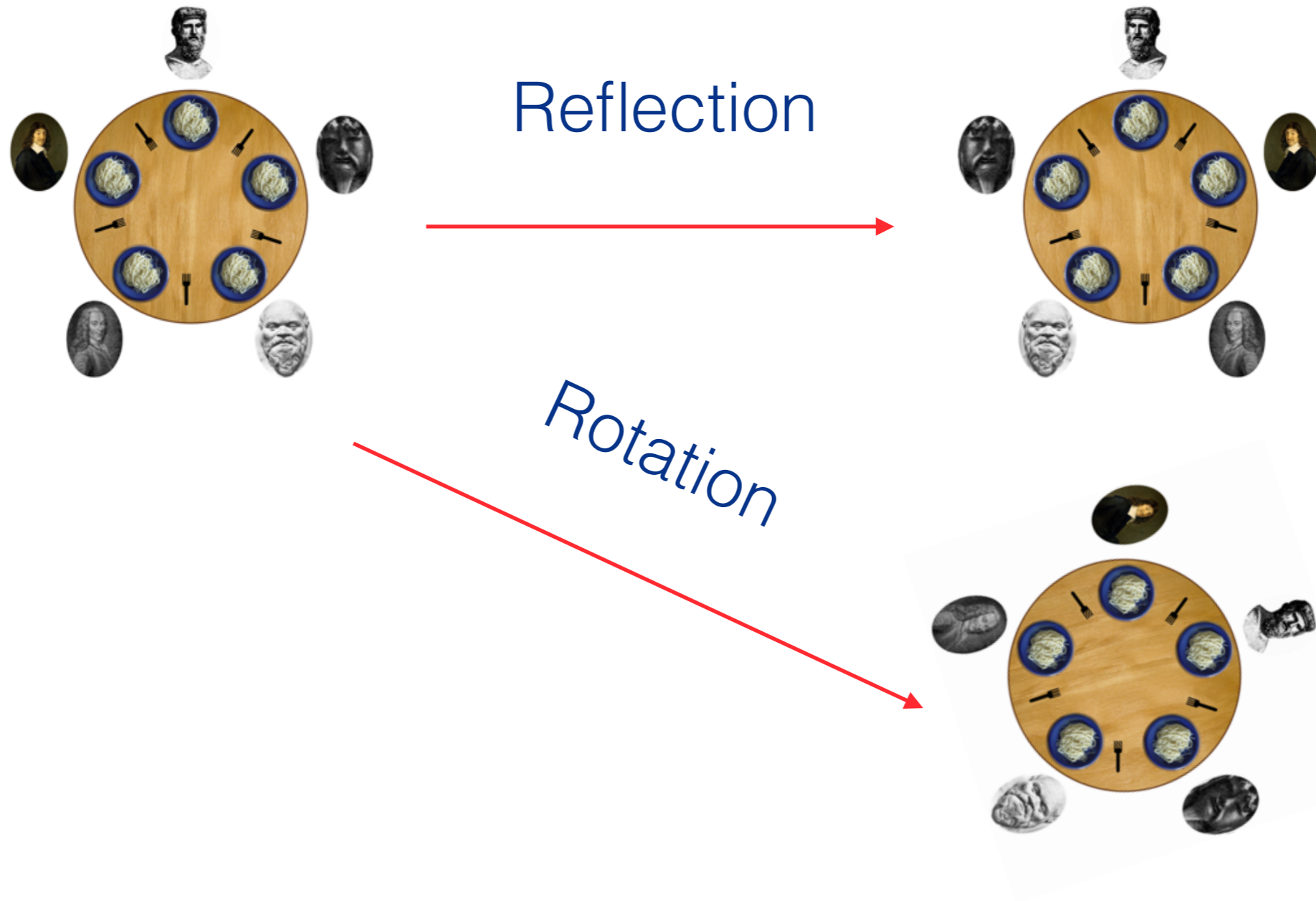
# Regular Symmetry Patterns

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# Symmetries in systems



# Symmetry examples



# Symmetries are closed under composition



# Symmetries as automorphisms

**Automorphism:** structure-preserving bijection on system configurations by permuting indices

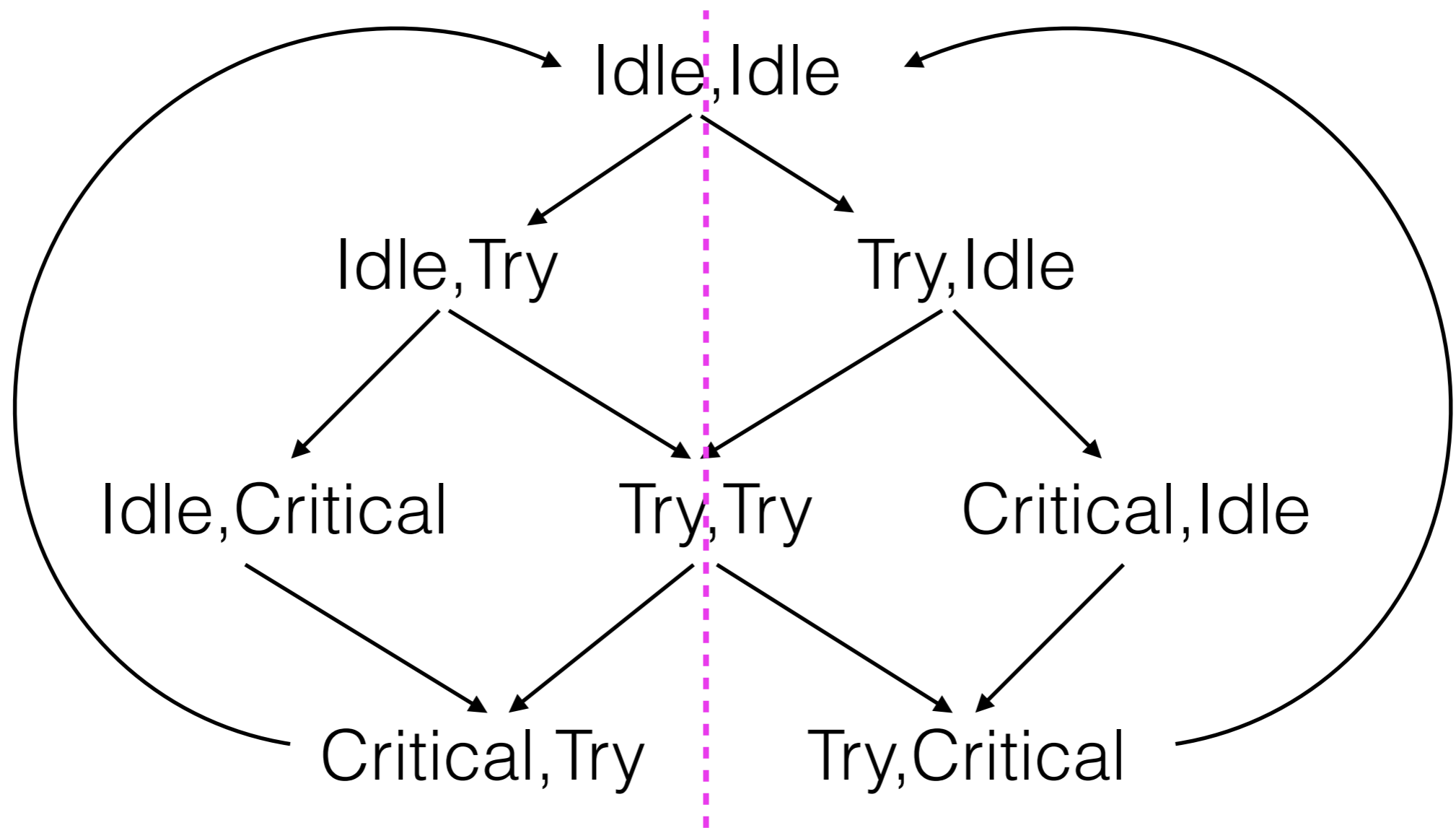
$$\pi : 1 \mapsto 2 \mapsto 3 \mapsto \dots \mapsto n \mapsto 1$$

(Critical)(Idle)(Idle)  $\longrightarrow$  (Idle)(Critical)(Idle)

The behaviour of systems is *indistinguishable* under an automorphism

---

# Automorphism example



Symmetry: 1  $\rightarrow$  2  $\rightarrow$  1



# Symmetries help model checking

**Gist:** Prune branches from states in the same equivalence class as visited states



*The space reduction can be exponential!*

*Works on all properties (safety, liveness, ...)*

# Two problems

Say, we mainly attack the first problem and, to some extent, the second problem.

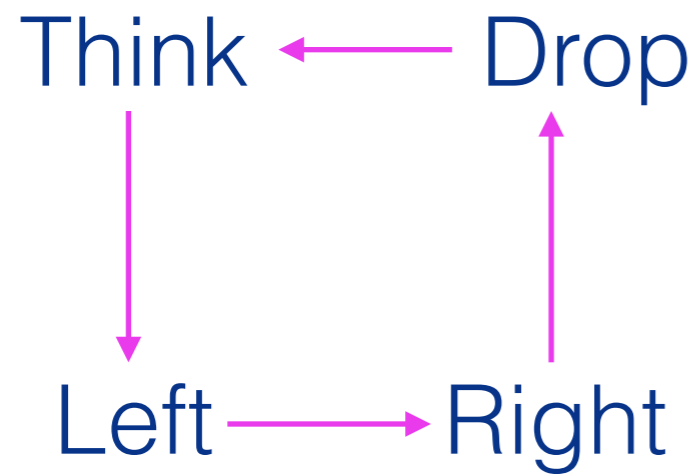
- **Symmetry identification:** how to identify symmetries in a given system
- **Symmetry exploitation:** (1) once symmetries are identified, check two states are similar (up to symmetries), (2) compute the “quotient” systems

*Both problems are in general computationally difficult!*

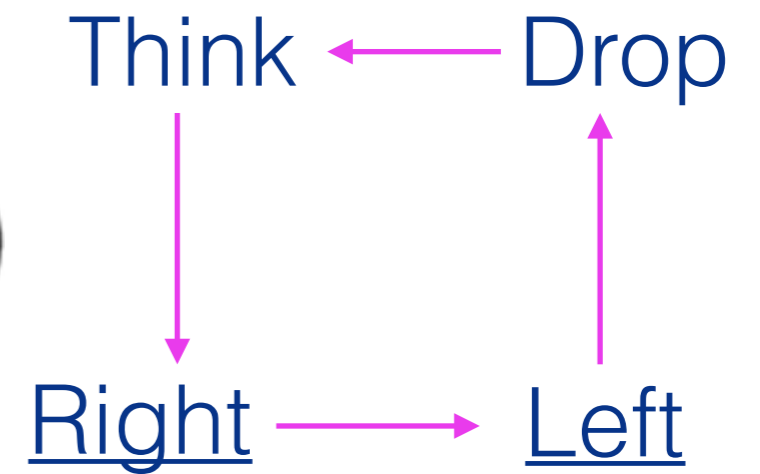
**Challenge: devise practical solutions to the problems**



# Concurrency by Replication



**Template 1**



**Template 2**

# Parameterised systems

Instance with *any* number of processes can be obtained by replicating templates (a.k.a. **parameterised systems**)

**Definition:** an *infinite* family of finite-state systems



# Parameterised Systems

## Help Verification

Instance-by-instance (using finite-state model checkers):

Size 1            0.1s

Size 2            0.1s

...

Size 5            1.5s

...

Size 10           62s

...

Size 15           Timeout

Parameterised verification (regular model checking, etc.):

Replication tends to produce “similar correctness proofs” for each size and can be symbolically represented

*Success on safety, but not so on other properties (e.g. liveness)*

# Can Parameterised Systems Help for Symmetry Finding?

Instance-by-instance (using finite-state symmetry finders):

Size 1	0.01s
Size 2	0.01s
...	
Size 5	0.2s
...	
Size 15	80s
...	
Size 20	Timeout

Parameterised:  
**??**

# Symmetry “Patterns” for Parameterised Systems

## **Observation:**

Instances of parameterised systems (obtained by) replications tend to exhibit *similar-looking symmetries*

# Pattern Example: Rotation



These 5 symmetries (case  $n=5$ ) can be generated by

$$\sigma_5 : 1 \mapsto 2 \mapsto 3 \mapsto 4 \mapsto 5 \mapsto 1$$

For general  $n$ , this rotation symmetry pattern is

$$\sigma_n : 1 \mapsto 2 \mapsto 3 \mapsto \dots \mapsto n \mapsto 1$$



# Pattern Example: Reflection



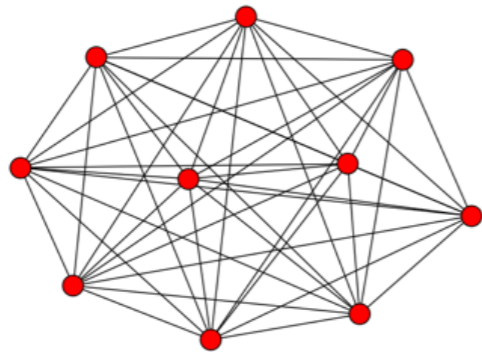
These 2 symmetries (case  $n=5$ ) can be generated by

$$\pi_5 : (1, 5)(2, 4)(3) \quad (\text{in cycle notation})$$

For general  $n$ , the reflection pattern is

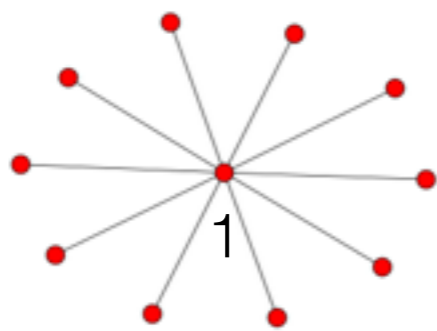
$$\pi_n : (1, n)(2, n-1) \cdots (\lfloor n/2 \rfloor, \lceil n/2 \rceil)$$

# Other patterns



Broadcast protocol

Full symmetry (all permutations on  $\{1, \dots, n\}$ )



Resource allocator

Full symmetry on subsystem (all permutations on  $\{1, \dots, n\}$  that fix the center point 1)

# Contributions

**Symbolic Framework** for Symmetry Patterns in Parameterised Systems

Language for Describing Systems: letter-to-letter transducers (standard in **regular model checking**)

**Language for Describing Symmetries: letter-to-letter transducers (NEW)**

**Expressive for describing practical symmetry patterns**

**automatic verification and synthesis of symmetry patterns**

# Symmetry verification

Does the given parameterised system exhibit ...?

- Rotations
- Reflections
- Full symmetries
- Above symmetries in a subsystem ...

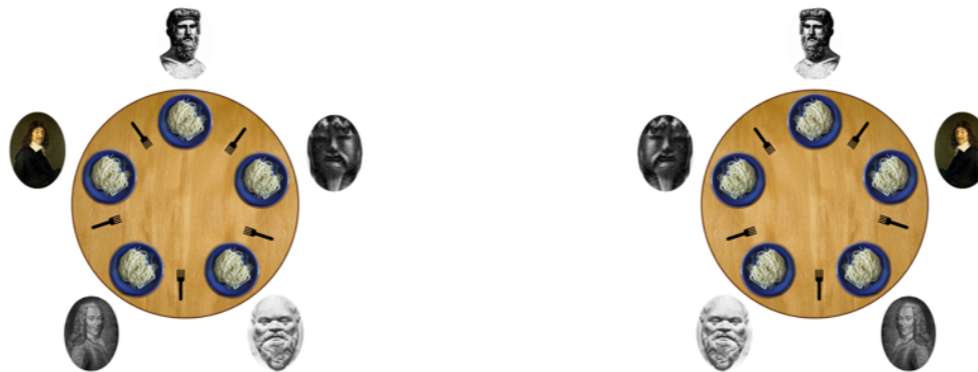
**Key Contribution:** *Each can be expressed and automatically checked in our framework!*

*Good news: there is a “library” of common symmetries*

# Symmetry synthesis

Symmetries in parameterised systems may not be obvious ...

- Data symmetries (e.g. fork position swapped)



- Symmetries in a subsystem (but which?)

**Contribution:** *a CEGAR method for synthesising symmetry patterns in a parameterised system*

The symbolic framework:  
more technical details



# Transducers

(Finite) Automata over the alphabet  $\Sigma \times \Sigma$

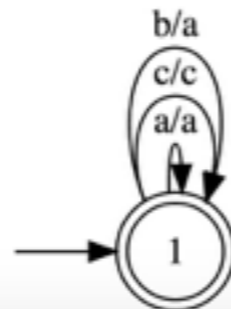
Symbolic representations of infinite binary relations

**Example:**

$$\Sigma = \{a, b, c\}$$

$$R = \{(v, w) : w \text{ is } v \text{ with } b \text{ replaced by } a\}$$

Automaton:



a	b	c
a	a	c

# Automatic transition systems (Regular Model Checking)

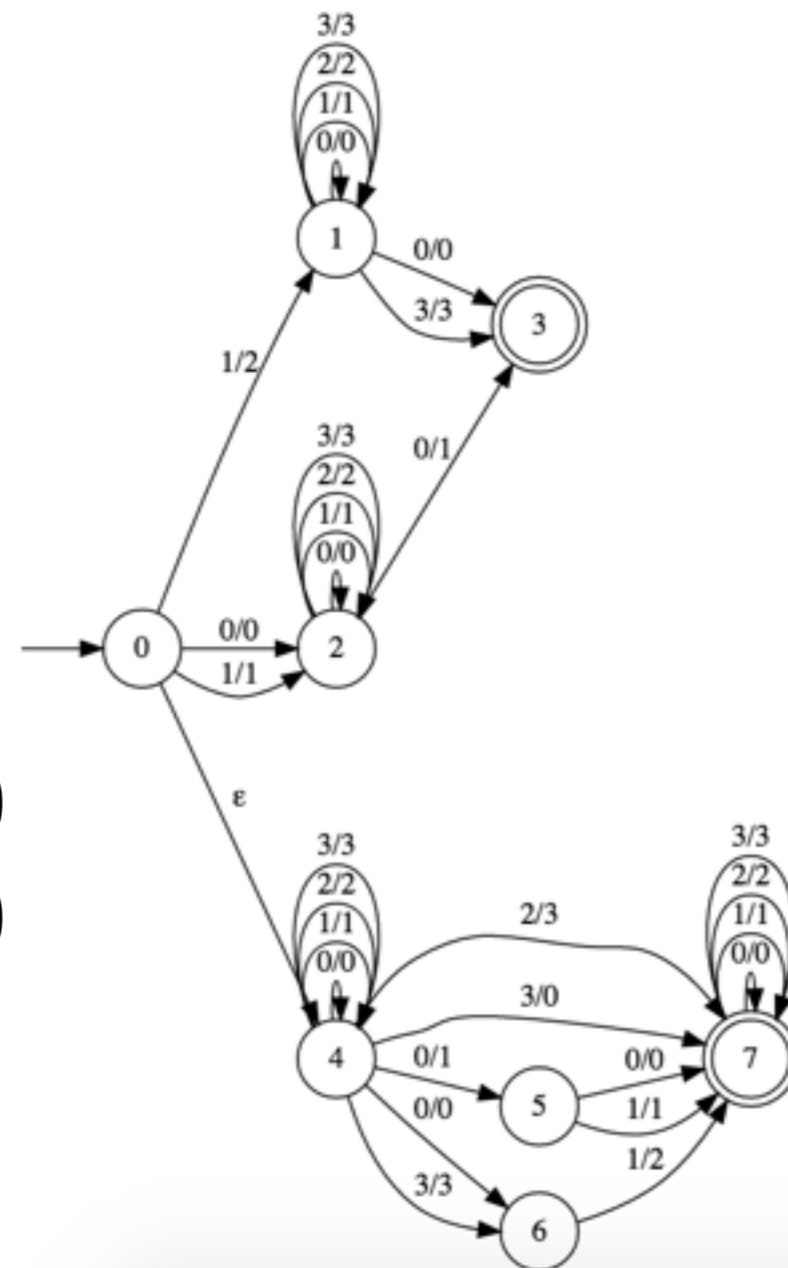
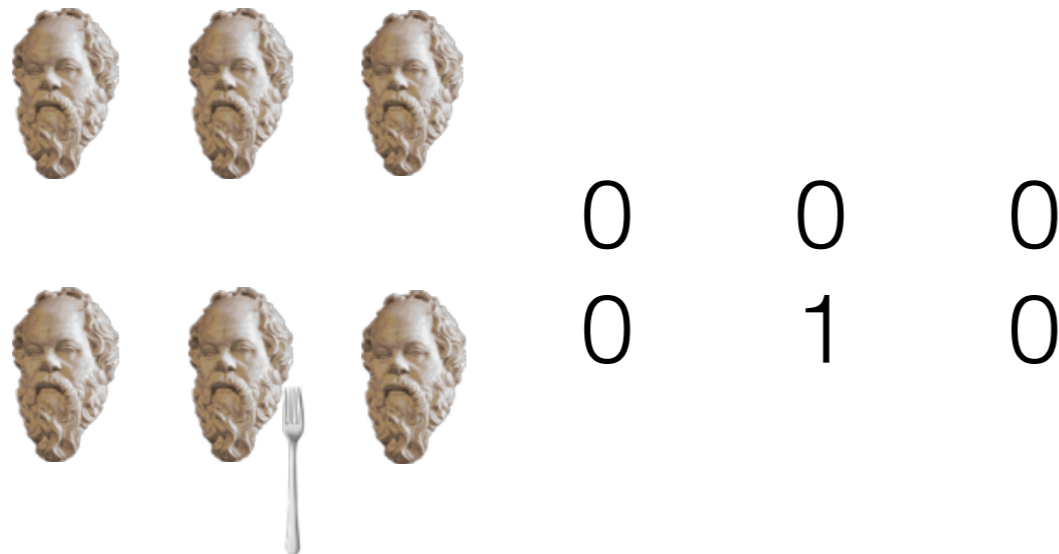
**Set of states:**  $\Sigma^*$  (or a regular subset thereof)

**Labelled transitions:** defined by a finite family of transducers (one transducer for each action label)

# Example: Dining-Philosopher (pick left first)

$$\Sigma = \{0, 1, 2, 3\}$$

0 - Thinking  $\longleftarrow$  3 - Drop Left  
 1 - Pick Left  $\longrightarrow$  2 - Pick Right



# Symmetry Pattern

Defn: a length-preserving *automorphism* on an automatic transition system

$$f : \Sigma^* \rightarrow \Sigma^*$$

$$\text{len}(f(v)) = \text{len}(v)$$

Bijection, Homomorphism, ...

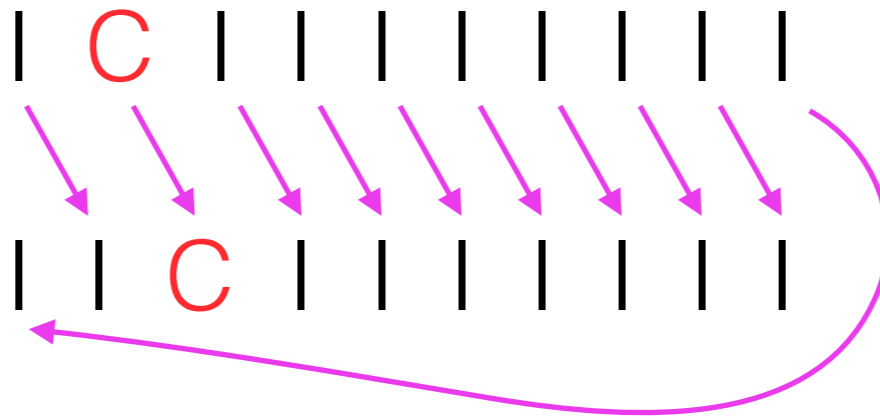
# Regular Symmetry Pattern

Defn: Symmetry pattern that can be represented by a transducer

View a function as a binary relation

**Examples (next few slides)**: rotation, swap, ...

# Rotation is regular



Automaton remembers when reading  $i$ th position:

1.  $i$ th position, 1st letter
2. 1st position, 2nd letter



# Symmetry Pattern Verification

# Verifying Regular Symmetry Patterns

**Theorem:** Checking whether a given automatic system exhibits a given regular symmetry pattern is PTIME checkable

Proof Idea: automata construction

**Corollary:** Checking whether a given automatic system exhibits a rotation symmetry is PTIME checkable

# Full Symmetry Pattern

All permutations on  $\{1, \dots, n\}$

This corresponds to  $n!$  automorphisms

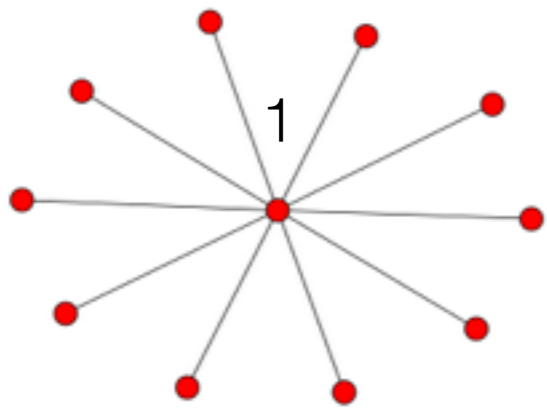
**Key:** the set of automorphisms forms a group under functional composition generated by:

$(1,2)$  — a swap

$(1, \dots, n)$  — a rotation

*Swap is also regular!*

# Full Symmetry in a Subsystem



All permutations on  $\{1, \dots, n\}$  that fix 1

This corresponds to  $(n-1)!$  automorphisms

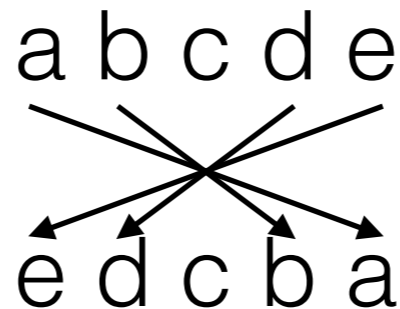
*These can be generated by  $(2,3)$  and  $(2,3,\dots,n)$*

# Verifying full symmetry

**Corollary:** Checking whether a given automatic system exhibits a full symmetry pattern (in a fixed subsystem) is PTIME checkable

# What about reflection?

Unfortunately, it is NOT regular!



You have to compare the first half of the string  
with the second half of the string

# Verifying reflection symmetry

**Theorem:** Checking whether a given automatic system exhibits a given reflection symmetry pattern is PTIME checkable

Proof idea: introduce a subclass of pushdown automata called

*Height-Unambiguous Pushdown Automata*

Key Property: they can be synchronised (unlike general PDA)

*Automatic symmetry verification extends to huCF patterns*

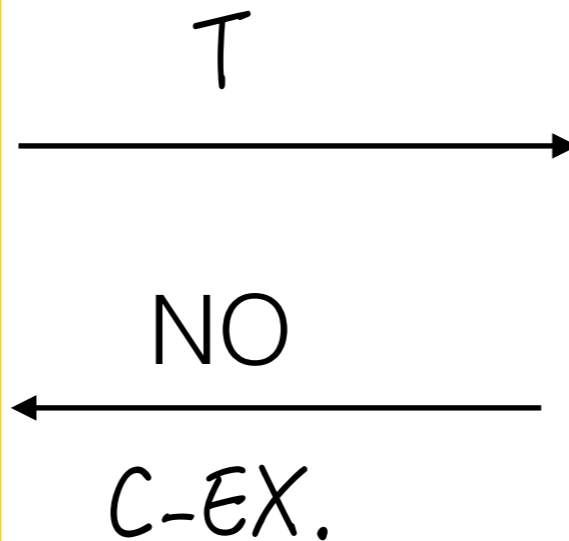
# Symmetry Pattern Synthesis



# Synthesise-Verify Loop

## Synthesise (SAT-solver)

1. Encode Transducers  $\mathcal{T}$  as Boolean Formulas
2. Maintain a set  $\mathcal{S}$  of boolean constraints that  $\mathcal{T}$  has to satisfy
3. Initialise  $\mathcal{S}$  to constraints like  $\mathcal{T}$  is not trivial,  $\mathcal{T}$  is infinite, ...



## Verify (automata method)

1. Is  $\mathcal{T}$  a (partial) function?
2. Is  $\mathcal{T}$  total?
3. Is  $\mathcal{T}$  injective?
4. Is  $\mathcal{T}$  surjective?
5. Is  $\mathcal{T}$  a homomorphism?

YES

FINISH

“Smart” enumeration of regular symmetry patterns:  
guess a transducer with 1 state, 2 states, 3 states,  
4 states, ...

# Counterexamples

Three forms of counterexamples:

1.  $v$  has to be included in the domain of  $T$
2.  $w$  has to be included in the range of  $T$
3. One of two contradictory pairs  $T(v,w)$  and  $T(v',w')$  must be eliminated.

*Each can be encoded as a boolean constraint!*

# Synthesis of Finite Existential Abstractions (for Proving Safety)

## Verify (automata method)

1. Is  $\tau$  a (partial) function?
2. Is  $\tau$  total?
3. Is  $\tau$  injective?
4. Is  $\tau$  surjective?
5. Is  $\tau$  a homomorphism?

Relax (3) and (4) in our synthesis-verify loop

Add to Synthesis (boolean constraint):  
- “The range of  $\tau$  finite?”

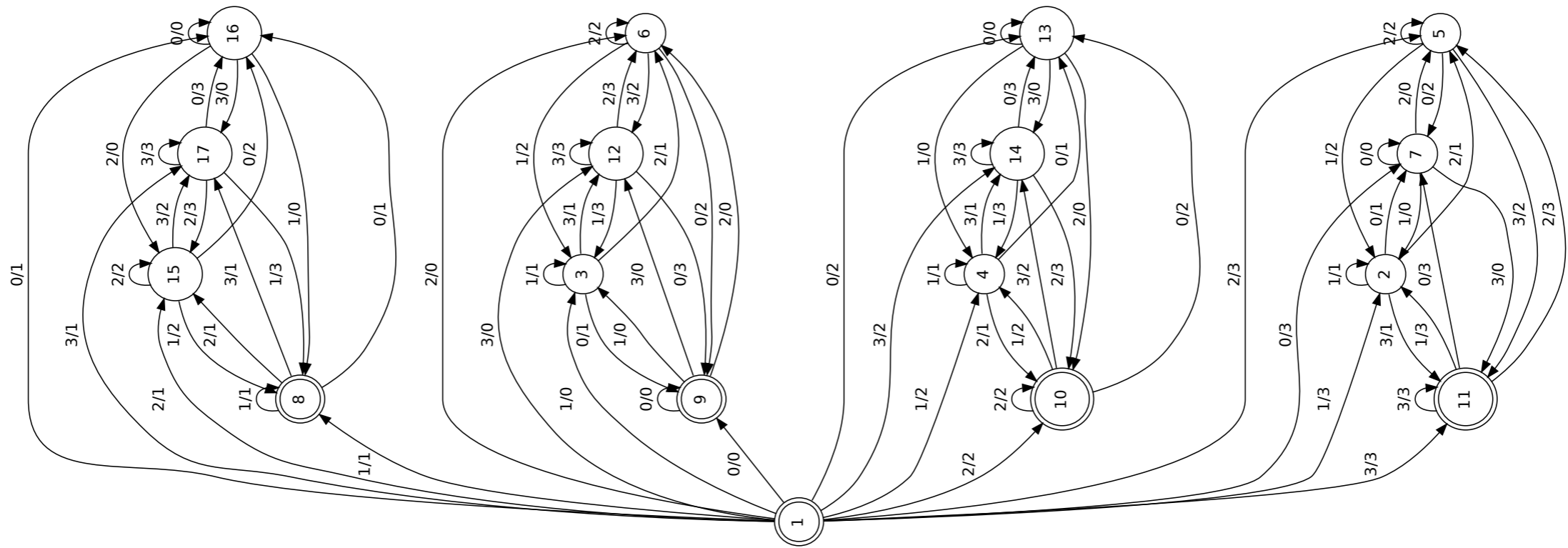
Add to Verify:  
- “Does the abstraction satisfy safety?”

*Can automatically check safety with a simple fixpoint computation (will terminate since range of  $\tau$  is finite)*

# Experiments and Examples

<b>Symmetry Systems (#letters)</b>	<b># Transducer states</b>	<b>Verif. time</b>	<b>Synth. time</b>
Herman Protocol (2)	5	0.0s	4s
Israeli-Jalfon Protocol (2)	5	0.0s	5s
Gries's Coffee Can (4)	8	0.1s	3m19s
Resource Allocator (3)	11	0.0s	4m56s
Dining Philosopher (4)	17	0.4s	26m

# Synthesised Transducer for Dining Philosopher



# Conclusion and Future Work

# Conclusion

- Look for symmetry patterns instead of symmetries (for an individual instance)
- Expressive symbolic framework for automatically verifying and synthesising symmetry patterns

# Future Work

- Synthesis of huCF symmetry patterns
- Synthesis of multiple symmetry patterns